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On the state of the art in proactive/reactive project scheduling

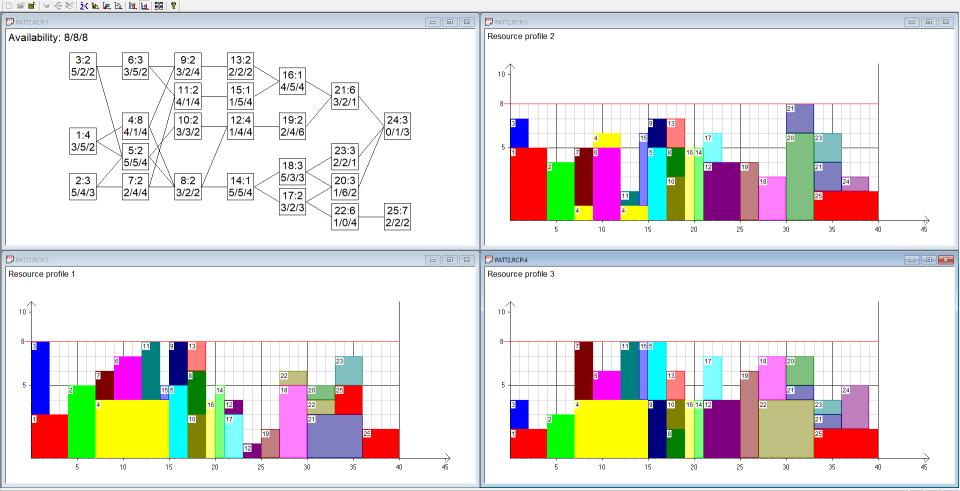
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https://feb.kuleuven.be/rescon/

RESCON - PAT72.RCP

File Appearance Algorithm Window Help



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Idea of robust project scheduling

- Generate a baseline schedule that incorporates a degree of anticipation of variability during project execution and/or information about the reactive scheduling approach to be used
- Objectives:
 - Solution robustness (stability):
 - Measure of the difference between the baseline schedule and the realized schedule
 - Quality robustness:
 - Sensitivity of the schedule performance in terms of the objective value (makespan) other than stability

The stochastic RCPSP

- Classes of scheduling policies:
 - Resource-based policies
 - Early-start policies
 - Preselective policies
 - Linear preselective policies
 - Activity-based policies
 - Pre-processing policies
- Drawback: no baseline schedule, thus no measure of solution robustness



Proactive-reactive project scheduling





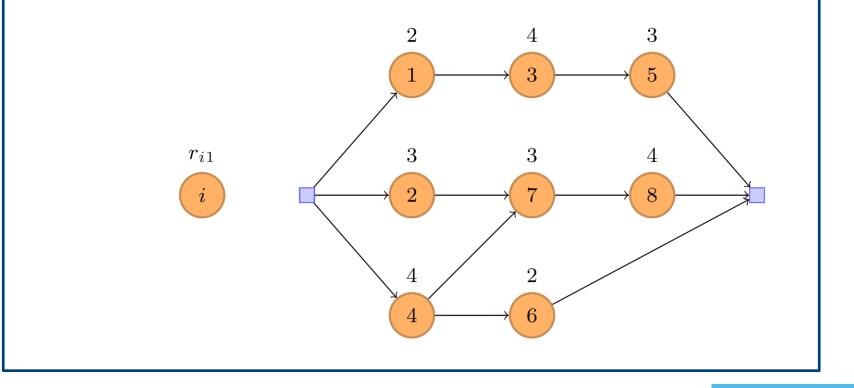
- Proactive scheduling:
 - Construct a robust baseline schedule that accounts for the available statistical knowledge of uncertainty and that is protected as best as possible against disruptions
- Reactive scheduling:
 - Revise or reoptimize a schedule whenever a schedule breakage occurs



The proactive and reactive RCPSP

An instance of the RCPSP

This instance contains <u>eight</u> activities and two dummy activities. This instance has <u>one</u> resource type of availability 8.



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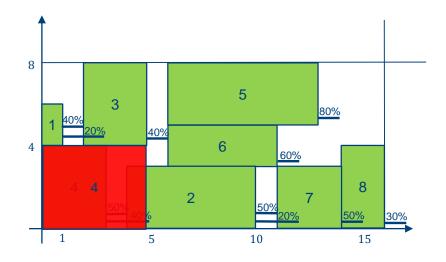
Example for \tilde{p}

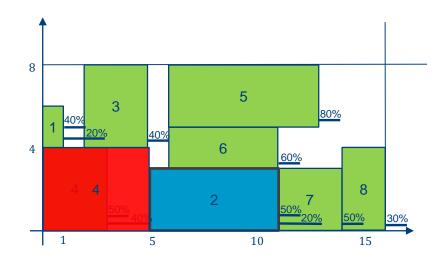
	$\pi(\tilde{p}_i = p)$									$w_{i,0}$	
	p = 0	p = 1	p=2	p = 3	p = 4	p = 5	p = 6	p = 7	p = 8		
\tilde{p}_0	1	-	-	-	-	-	-	-	-	-	Total number of
$ ilde{p}_1$	-	0.4	0.4	0.2	-	-	-	-	-	4	combinations:
\tilde{p}_2	-	-	-	-	-	-	0.3	0.5	0.2	4	
$ ilde{p}_3$	-	-	-	0.6	0.4	-	-	-	-	7	$2^5 \times 3^3 = 864$
\tilde{p}_4	-	-	-	0.1	0.5	0.4	-	-	-	1	
\tilde{p}_5	-	-	-	-	-	-	-	0.2	0.8	4	
\tilde{p}_6	-	-	-	-	-	0.4	0.6	-	-	1	
\tilde{p}_7	-	-	-	0.5	0.5	-	-	-	-	1	
\tilde{p}_8	_	-	0.7	0.3	-	-	-	-	-	1	
\tilde{p}_9	1	-	-	-	-	-	-	-	-	38	

An example realization

 $\mathbf{p_1} = (0,2,8,3,5,7,5,4,2,0) \\ \pi(\widetilde{\mathbf{p}} = \mathbf{p_1}) = 0.054\%$

Proactive and reactive scheduling



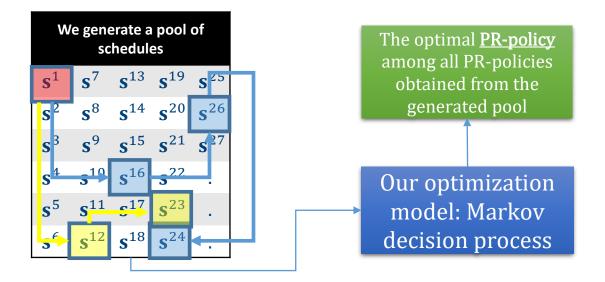


Proactive solution

Reactive scheduling



The basic idea



Objective: to minimize the expected value of (cost of the baseline schedule + cost of a series of reactions)

A PR-policy

PR-pc	blicy Π for realization	\mathbf{p}_l			
t = 0		$t = t_1$	<i>t</i> =	t ₂ t =	= t ₃
	$s^{[0]}\pi_{l}$ (baseline schedule)	<u>,</u>	_S [1] _{П,l} П	$S^{[2]_{\Pi,l}}$	¹ _S ^[3] Π, <i>ι</i>
	Disruption 1	for $S^{[0]_{\Pi,l}}$ at time t_1	Disruption for $S^{[1]_{\Pi,l}}$ at t	time t_2 Disru	ption for $S^{[2]_{\Pi,l}}$ at time t_3

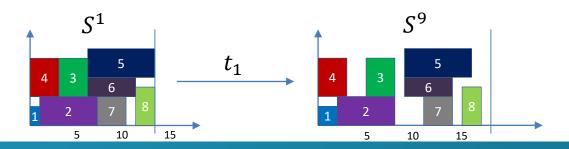
A chain of reactions dictated by PR-policy
$$\Pi$$
 for \mathbf{p}_l
 $CH_{\Pi,l}: S^{[0]}_{\Pi,l} \xrightarrow{t_1} S^{[1]}_{\Pi,l} \xrightarrow{t_2} \dots \xrightarrow{t_{\nu_{\Pi,l}}} S^{[\nu_{\Pi,l}]}_{\Pi,l} \xrightarrow{\text{The number of reactions for the combination }(\Pi, l)}$

$$\Pi = \{CH_{\Pi,1}, \dots, CH_{\Pi,|\mathbf{p}|}\}$$

An example

PR-policy Π_1

									•	
$\mathbf{p_1} = (0,2,8,3,5,7,5,4,2,0)$	S^1	S^2	S^3	S^4	S^5	S^6	S^7	S^8	S^9	S^{10}
$CH_{\Pi_1,1}: S^1 \xrightarrow{t_1=1} S^9$	0	0	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0	0	0
	1	1	0	1	5	0	7	4	2	7
	3	3	4	4	3	3	3	3	5	5
	0	0	4	0	0	7	0	0	0	9
	6	6	7	7	7	7	7	7	9	14
	6	6	7	7	7	12	5	7	9	14
	7	8	7	8	12	12	14	12	11	15
$\mathbf{\Pi} = (C \Pi - C \Pi)$	11	13	13	12	15	15	17	15	15	20
$\mathbf{\Pi_1} = \{CH_{\Pi_1,1}, \dots, CH_{\Pi_1, \mathbf{p} }\}$	13	15	15	15	17	18	19	18	18	23



Deadchains

A chain of reactions dictated by PR-policy Π for \mathbf{p}_l

$$CH_{\Pi,l}: S^{[0]_{\Pi,l}} \xrightarrow{t_1} S^{[1]_{\Pi,l}} \xrightarrow{t_2} \dots \xrightarrow{t_{\nu_{\Pi,l}}} S^{[\nu_{\Pi,l}]_{\Pi,l}}$$

What if this schedule is not feasible for \mathbf{p}_l ?

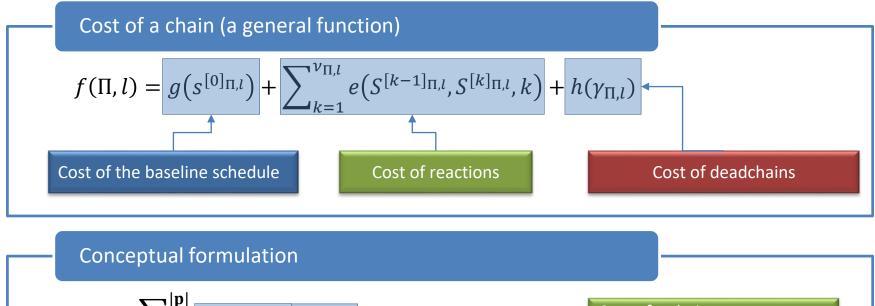
What is a deadchain?

 $S^{[\nu_{\Pi,l}]_{\Pi,l}}$ is a deadend Ω if it is <u>not</u> feasible for realization \mathbf{p}_l . $CH_{\Pi,l}$ is a deadchain if it contains a deadend.

$\gamma_{\Pi,l} \in \{0,1\}$

 $\gamma_{\Pi,l} = 1$ if chain $CH_{\Pi,l}$ is a deadchain $\gamma_{\Pi,l} = 0$ otherwise

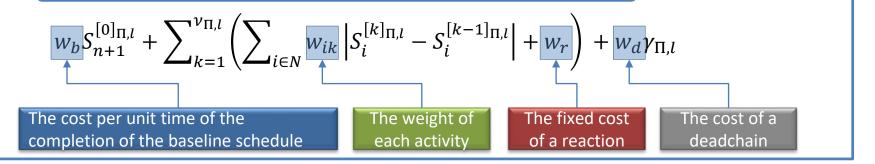




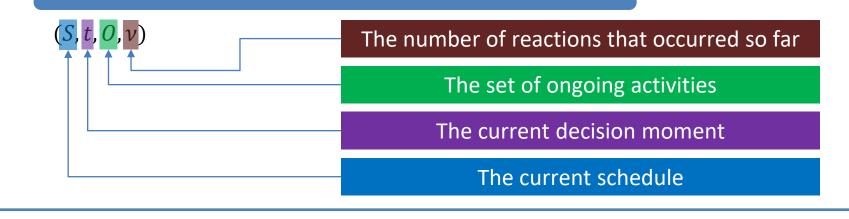
P:
$$\min_{\Pi \in \Xi} \sum_{l=1}^{|\mathbf{p}|} \pi(\tilde{\mathbf{p}} = \mathbf{p}_l) f(\Pi, l)$$

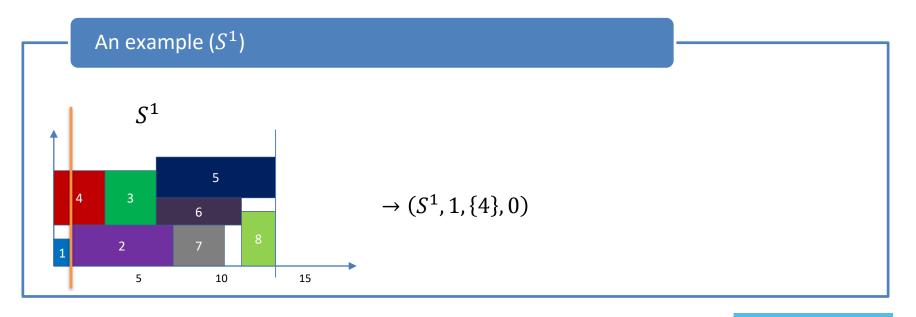
Probability of occurrence

Cost of a chain (an example: used in our method)



Representation of states in Model 1





A chance transition

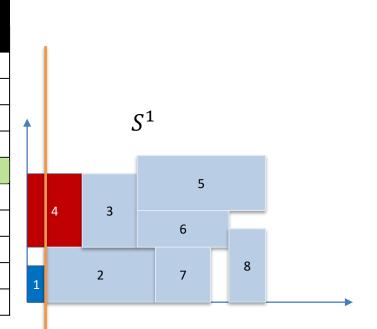
$$(S^1, 0, \emptyset, 0)$$

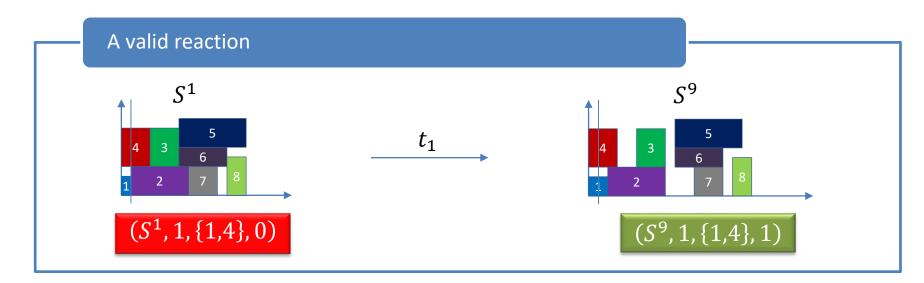
0.40
 $(S^1, 1, \{4\}, 0)$
0.60
 $(S^1, 1, \{1,4\}, 0)$

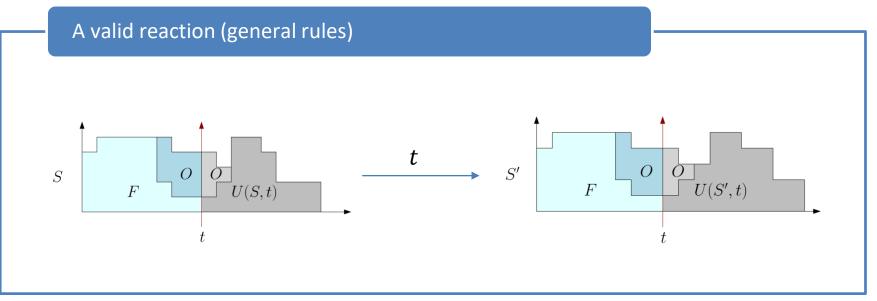
Continue with the same schedule

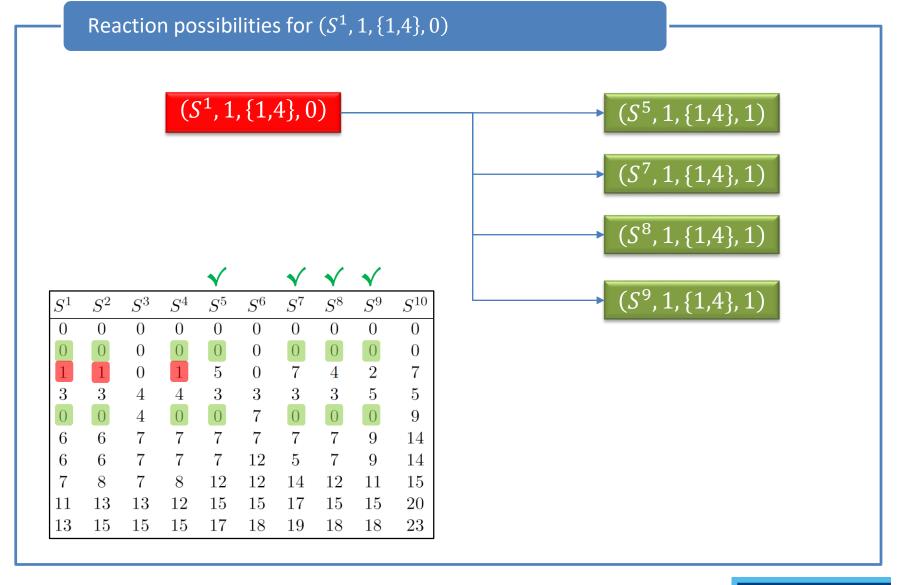
Infeasible: needs reaction

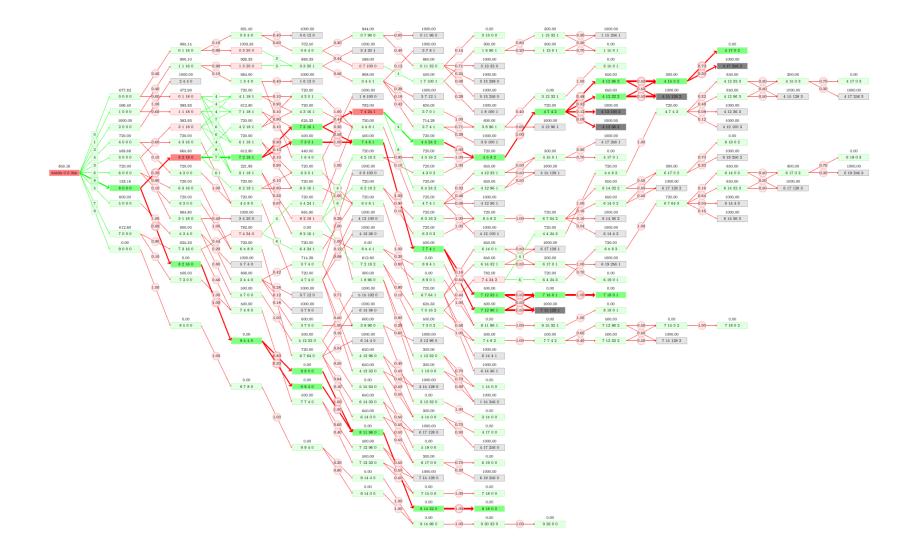
	$\pi(ilde{p}_{ ext{i}}=p)$										
	0	1	2	3	4	5	6	7	8		
${\widetilde p}_0$	1	-	-	-	-	-	-	-	-		
\widetilde{p}_1	-	0.4	0.4	0.2	-	-	-	-	-		
\widetilde{p}_2	-	-	-	-	-	-	0.3	0.5	0.2		
\widetilde{p}_3	-	-	-	0.6	0.4	-	-	-	-		
\widetilde{p}_4	-	-	-	0.1	0.5	0.4	-	-	-		
${\widetilde p}_5$	-	-	-	-	-	-	-	0.2	0.8		
${\widetilde p}_6$	-	-	-	-	-	0.4	0.6	-	-		
\widetilde{p}_7	-	-	-	0.5	0.5	-	-	-	-		
${\widetilde p}_8$	-	-	0.7	0.3	-	-	-	-	-		
\widetilde{p}_9	1	-	-	-	-	-	-	-	-		

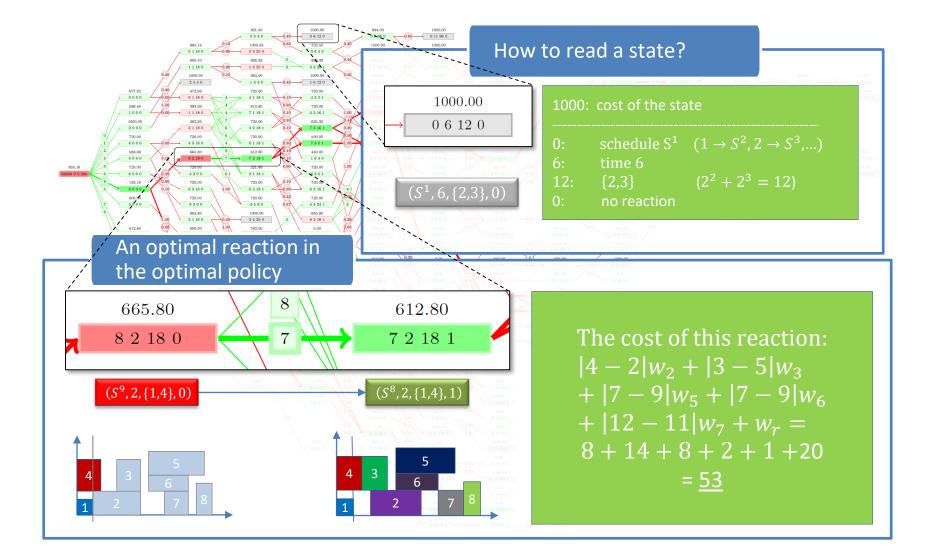


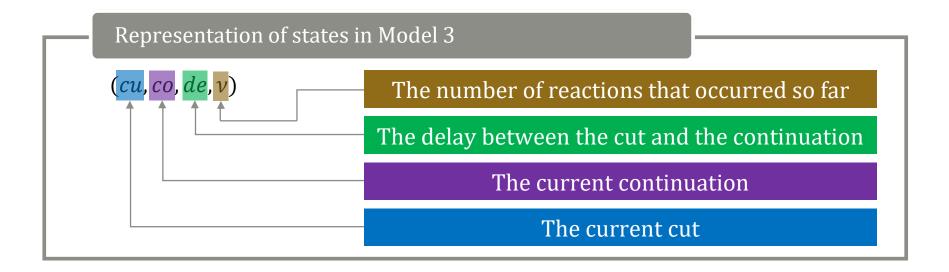




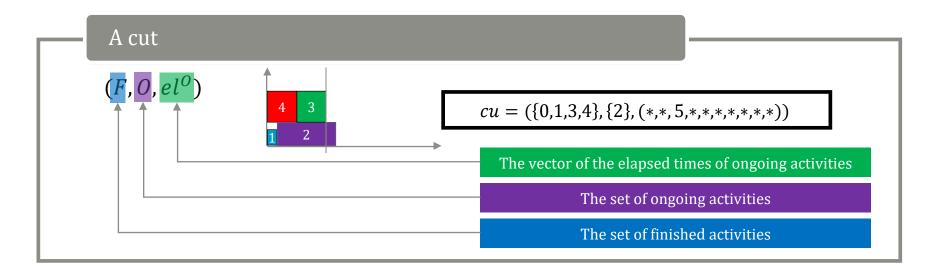


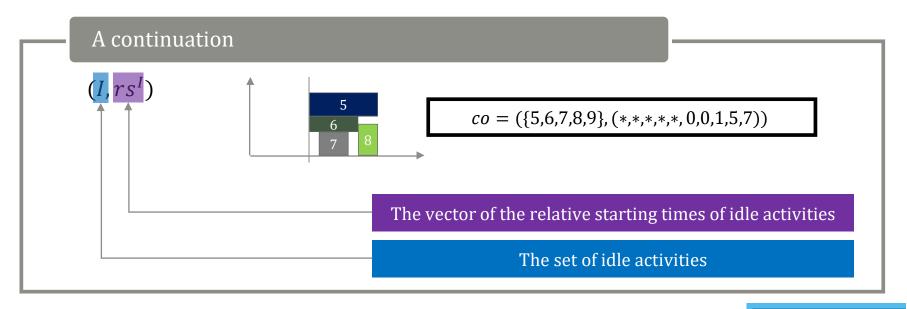


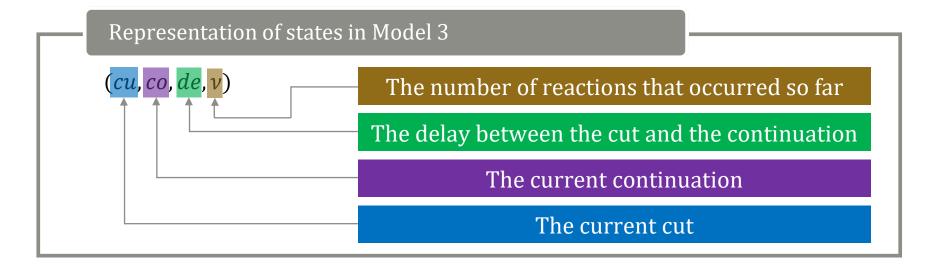


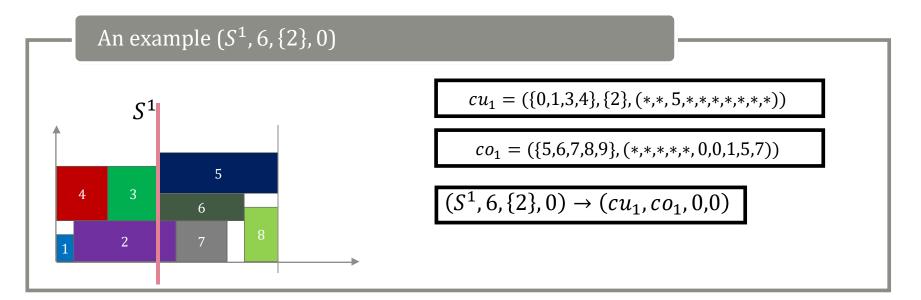






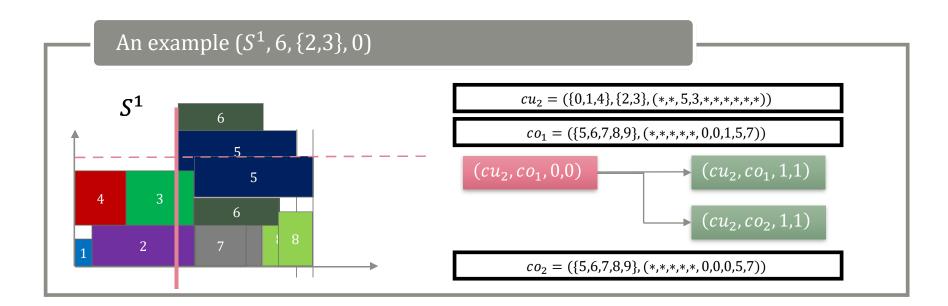




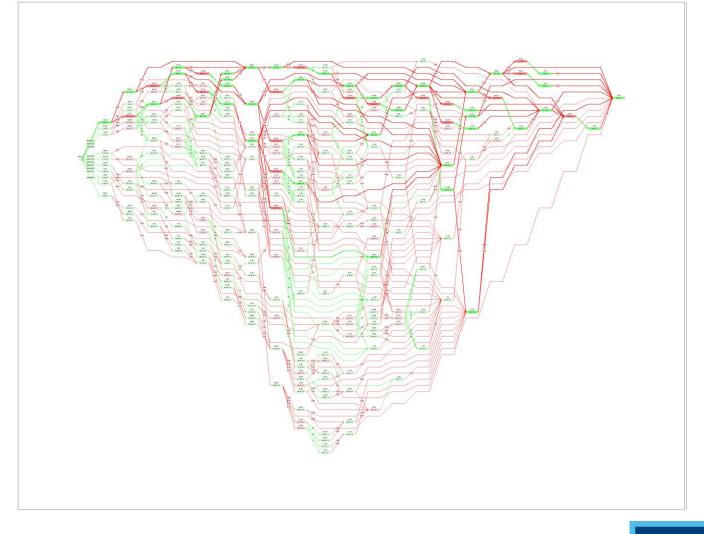


State space

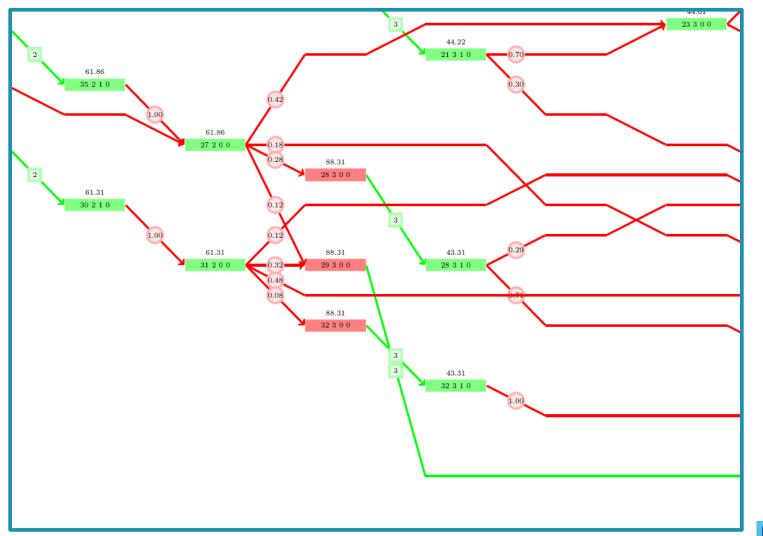
- We first generate all possible cuts that can be obtained from **S**
- We generate all possible continuations that can be obtained from **S**
- $de \in \{0,1\}$
- There is a state for each combination (cu, co, de, v) such that $(F \cup O) \cap I = \emptyset$ and $(F \cup O) \cup I = N$



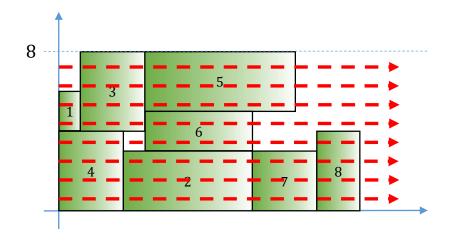
Model 3: Markov decision process



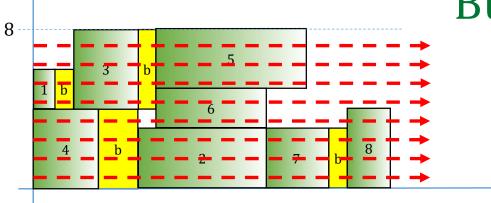
Model 3: optimal PR-policy



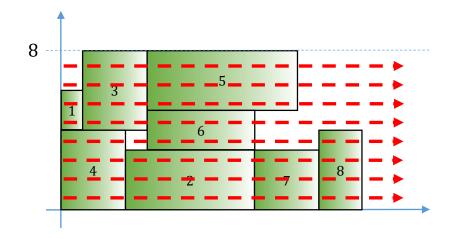
Selection- and buffer-based reactions



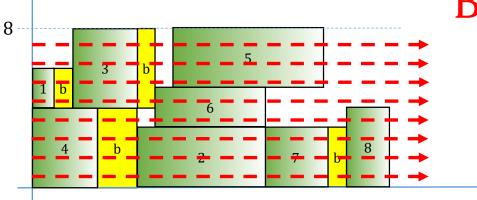
Selection-based ✓ Buffer-based ✓



A selection- but not buffer-based reaction



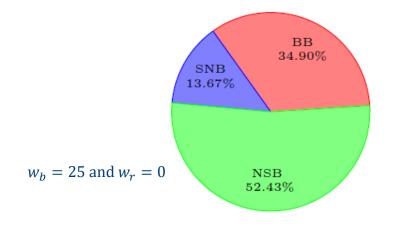
Selection-based ✓ Buffer-based X

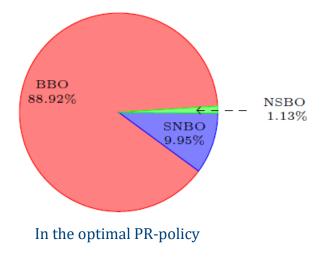


Computational results

Three exclusive classes of reactions

The class of non-selection-based (NSB) reactions The class of selection but not buffer-based (SNB) reactions The class of buffer-based (BB) reactions





Selection-based

References

- Davari, M., 2017, Contributions to complex project and machine scheduling problems, PhD dissertation Faculty of Economics and Business, KU Leuven, number 554.
- Davari, M. & E. Demeulemeester, 2019, The proactive and reactive resourceconstrained project scheduling problem, Journal of Scheduling, 22 (2), 211-237.
- Davari, M. & E. Demeulemeester, 2019, Important classes of reactions for the proactive and reactive resourceconstrained project scheduling problem, Annals of Operations Research, 274 (1-2), 187-210.



